

# Repetition Administration of Multipath Routing for Intrusion Tolerance in Heterogeneous Wireless Sensor Networks



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**Abstract:** This project proposed redundancy management of heterogeneous wireless sensor networks (HWSNs), to manage efficient communication in heterogeneous wireless sensor network by utilizing multipath routing to respond the user queries. The main aim of this project is to exploit the trade off between energy consumption vs. reliability, timeliness, and security to maximize the system useful lifetime. The other problem in the HWSN is optimization in this project we trade of optimization problem by determining the best redundancy level to apply to multipath routing for intrusion tolerance. More ever we adopt voting-based distributed intrusion detection algorithm is applied to identify malicious nodes in typical HWSN. After that we implement a novel probability model to analyze the best redundancy level in terms of path redundancy and source redundancy, as well as the best intrusion detection settings in terms of the number of voters and the intrusion invocation interval under which the lifetime of a HWSN is maximized.

**Key Words:** *Dynamic keys, Distributed denial of service attacks, firewall, IP address spoofing, packet filtering.*

## 1. INTRODUCTION

Wireless sensor network (WSN) refers to a group of spatially dispersed and dedicated sensors for monitoring and recording the physical conditions of the environment and organizing the collected data at a central location. WSNs measure environmental conditions like temperature, sound, pollution levels, humidity, wind speed and direction, pressure, etc. The tradeoff issue between vitality utilization vs. Qos increase gets to be a great deal more confounded when inside assailants are available as a way may be broken when a noxious hub is on the way. This is particularly the case in heterogeneous

WSN (HWSN) situations in which CH hubs may take a more basic part in social occasion and directing sensing information. Therefore, likely the framework would utilize an interruption identification framework (IDS) with the objective to identify and uproot malevolent hubs. While the writing is rich in interruption recognition strategies for wsn [7]–[11], the issue of how frequently interruption identification ought to be conjured for vitality reasons so as to uproot possibly malevolent hubs so that the framework lifetime is expanded (say to keep a Byzantine disappointment [12]) is to a great extent unexplored. The issue is particularly discriminating for energyconstrained

WSNs intended to stay alive for a long mission time Multipath steering is viewed as a successful instrument

for shortcoming and interruption tolerance to enhance information conveyance in Wsn. The fundamental thought is that the likelihood of no less than one way arriving at the sink

hub or base station builds as we have more ways doing information delivery. while most earlier research concentrated on utilizing multipath directing to enhance unwavering quality [2], [3], [13], some consideration has been paid to utilizing multipath directing to endure insider assaults [14]. These studies, in any case, generally disregarded the tradeoff between Qospicup vs. vitality utilization which can antagonistically abbreviate the framework lifetime

The examination issue we are tending to in this paper is powerful excess administration of a bunched HWSN to delay its lifetime operation in the vicinity of temperamental furthermore malignant hubs. We address the tradeoff between vitality utilization vs. Qos pick up in unwavering quality, convenience and security with the objective to augment the lifetime of a grouped HWSN while fulfilling application Qos prerequisites in the setting of multipath directing

We consider this streamlining issue for the case in which a voting based disseminated interruption discovery calculation is connected to expel vindictive hubs from the HWSN. Our commitment is a model-based investigation philosophy by which the ideal multipath repetition levels and interruption discovery settings may be distinguished for fulfilling application Qos prerequisites while augmenting the lifetime of HWSN's.

## 2. Related Work

The WSN is built of "nodes" – from a few to several hundreds or even thousands, where each node is connected to one (or sometimes several) sensors. Each such sensor network node has typically several parts: a radio transceiver with an internal antenna or connection to an external antenna, a microcontroller, an electronic circuit for interfacing with the sensors and an energy source, usually a battery or an embedded form of energy harvesting. A sensor node might vary in size from that of a shoebox down to the size of a grain of dust, although functioning "motes" of genuine microscopic dimensions have yet to be created. The cost of sensor nodes is similarly variable, ranging from a few to hundreds of dollars, depending on the complexity of the individual sensor nodes.

Hierarchical or cluster-based routing, originally proposed in wireless networks, are well-known techniques with special advantages related to scalability and efficient communication. As such, the concept of hierarchical routing is also utilized to perform energy-efficient routing in WSNs. In a hierarchical architecture, higher energy nodes can be used to process and send the information while low energy nodes can be used to perform the sensing in the proximity of the target. This means

that creation of clusters and assigning special tasks to cluster heads can greatly contribute to overall system scalability, lifetime, and energy efficiency. Hierarchical routing is an efficient way to lower energy consumption within a cluster and by performing data aggregation and fusion in order to decrease the number of transmitted messages to the BS.

In case of WSNs, the definition of network lifetime is application specific [4]. It may be taken as the time from inception to the time when the network becomes non functional. A network may become non-functional when a single node dies or when a particular percentage of nodes perishes depending on requirement. Moreover it is accepted universally that balancing the energy dissipation by the nodes of the network is a key factor for prolonging the lifetime [4]. Each sensor node is provided with transmit power control and omni-directional antenna and therefore can vary the area of its coverage [2, 5]. It has been established in [1] that communication requires significant amount of energy as compared to computations.

Hence it is of utmost importance that sensor nodes collaborate with each other in an energy-efficient manner for transmitting and receiving data so that lifetime enhancement is achieved. In this paper, we consider a wireless sensor network where the base station is fixed and located far off from the sensed area. Furthermore all the nodes are static, homogenous and energy constrained and capable of communicating with the BS. The network being homogenous no high energy nodes are available hence communication between the nodes and the base station is expensive affair. Moreover all nodes have information about their respective distances from the BS in the static environment as stated in [2]. It has been often observed that sensor networks are burdened with redundant data during the process of systematic data gathering from the field into neighbours and strangers

In the connection of secure multipath directing for interruption tolerance, [2] gives an astounding overview in this subject. In [15] the creators considered a multipath directing convention to endure dark gap and specific sending assaults. The essential thought is to utilize catching to abstain from sending bundles to malevolent hubs. In [14] the creators considered a disjoint multipath directing convention to endure interruption utilizing various disjoint ways in WSNs. Our work additionally utilizes multipath directing to endure interruption. Notwithstanding, we particularly consider vitality being expended for interruption identification, and both CHs and Sns could be traded off for lifetime augmentation. In [3] a randomized dispersive multipath steering convention is proposed to maintain a strategic distance from dark gaps

### 3. Probability Model

In this area we create a likelihood model to gauge the MTTF of a HWSN utilizing multipath information sending to reply inquiries issued from a portable client wandering in the HWSN region. Table I gives the documentation used to images and their physical meanings. We utilize the same documentation for both CHs and Sns, e.g., Pfp and Pfn. At the point when separating a CH from a SN is fundamental, we utilize the superscripts or subscripts "CH" and "SN", e.g.,

PCH fp and PCH fn for a CH, and PSN fp and PSN fn for a SN. A parameter is named as data, inferred, plan or yield. Specifically, mp (way repetition), ms (source excess), m (the amount of voters for interruption location) and TIDS (the interruption location interim) are outline parameters whose qualities are to be distinguished to amplify MTTF, when given a set of info parameter qualities characterizing the operational furthermore natural conditions. Inferred parameters are those determining from data parameters. There stands out yield parameter, specifically, MTTF. Note that most inferred parameters are dynamic, i.e., as a capacity of time. For instance, hub thickness meant by  $\lambda(t)$  diminishes about whether as a result of hub disappointment/ousting as time advances

### 4. System Dynamics

At first at arrangement time all hubs (CHs or Sns) are great hubs. Accept that the catch time of a SN takes after a circulation capacity  $F_c(t)$  which could be resolved based on authentic information and learning about the target application nature. At that point, the likelihood that a SN is traded off at time  $t$ , given that it was a decent hub at time  $t - TIDS$ , signified by  $P_c$ , is given by:

$$P_c = 1 - P \{x > t | x > t - Tids\} = 1 - P \{x > t, x > t - Tids\} \\ P \{x > t - Tids\} = 1 - 1 - F_c(t) 1 - F_c(t - TIDS) \quad (2)$$

We note that  $P_c$  is time subordinate. For the exceptional case in which the catch time is exponential conveyed with rate  $\lambda_c$ ,  $P_c = 1 - e^{-\lambda_c \times tids}$ . Review that the voting-based conveyed IDS executes intermittently with TIDS being the interim. At the  $i$ th IDS execution time (signified by  $t_{i,i}$ ), a great hub might have been traded off with likelihood  $P_c$  since the past IDS execution time ( $t_{i,i-1}$ ). Let  $ngood(t)$  and  $nbad(t)$  signify the amounts of great and terrible neighbor hubs at time  $t$ , individually, with  $ngood(t) + nbad(t) = n(t)$ . At that point, the populace of great and terrible neighbor hubs at time  $t_{i,i}$  just before IDS execution might be recursively assessed from the populace of great and terrible neighbor hubs at time  $t_{i,i-1}$  as takes after:  $ngood(t_{i,i}) = ngood(t_{i,i-1}) - ngood(t_{i,i-1}) \times P_c$   $nbad(t_{i,i}) = n$

The main term in Eq. (4) represents the case in which more than 1/2 of the voters chose from the focus on hub's neighbors are awful sensors who, as a consequence of performing reviling assaults, will dependably vote a decent hub as an awful hub to break the usefulness of the HWSN. Since more than 1/2 of the  $m$  voters vote no, the target hub (which is a decent hub) is diagnosed as an awful hub for this situation, bringing about a false positive.

### 5. Query Success Probability

We will utilize the documentation  $S_{nj}$  to allude to SN  $j$  and  $Ch_j$  to allude to CH  $j$ . There are three routes by which information sending from  $Ch_j$  to  $Ch_k$  could fizzle: (a) transmission speed infringement; (b) sensor/channel disappointments; and (c)  $Ch_j$  is bargained. The primary wellspring of disappointment, transmission speed infringement, represents question due date infringement. To know the disappointment likelihood because of transmission velocity infringement, we first infer the base jump by-bounce transmission speed needed to fulfill the inquiry due date  $Treq$ . Let  $dsn_{-ch}$  be the normal separation between a SN (chose to

report sensor readings) and its CH and dch-pc be the normal separation between the source CH and the PC tolerating the inquiry result. Given an inquiry due date Treq as include, an information parcel from a SN through its CH to the PC must achieve the PC inside Treq. Subsequently, the base jump by-bounce transmission rate indicated by Sreq is given by:  $Sreq = dsn-ch + dch-p$

## 6. Energy Consumption

Presently we appraise the measures of vitality used amid a inquiry interim  $[tq,i, tq,i+1]$ , an IDS interim  $[ti,i, ti,i+1]$ , and a grouping interim  $[tc,i, tc,i+1]$ , to gauge  $Nq$ , the most extreme number of inquiries the framework can conceivable handle before running into vitality weariness. To standardize vitality utilization over  $Nq$  inquiries, let  $\alpha$  be the proportion of the IDS execution rate to the inquiry entry rate and let  $\beta$  be the degree of the grouping rate to the inquiry landing rate so that  $\alpha nq$  furthermore  $\beta nq$  are the amounts of IDS cycles and grouping cycles, separately, before framework vitality depletion. At that point, we can gauge  $Nq$  by the way that the aggregate vitality expended because of interruption location, grouping and inquiry preparing is equivalent to the framework vitality as takes after:

$$E_{init} = \alpha nq \sum_{i=1}^{EIDS} (ti,i) + \beta nq \sum_{i=1}^{Eclustering} (tc,i) + Nq \sum_{i=1}^{Eq} (tq,i)$$

Beneath we layout how to compute EIDS (ti,i), Eclustering (tc,i) and Eq (tq,i). We first gauge the measure of vitality devoured by transmission and gathering over remote connection. The vitality used by a SN to transmit an scrambled information bundle of length nb bits over a separation r is evaluated as [1]:

$$E_t = nb (E_{elec} + E_{amp,r})$$

Here  $E_{elec}$  is the vitality disseminated to run the transmitter also recipient hardware,  $E_{amp}$  is the vitality utilized by the transmit enhancer, and r is the transmission radio reach. We utilize the current rch and rsn to infer ECH T and ESN T. We set  $E_{amp} = 10$  pj/bit/m<sup>2</sup> and  $x = 2$  when  $d < d_0$ , and  $E_{amp} = 0.0013$  pj/bit/m<sup>4</sup> and  $x = 4$  overall. The vitality used by a hub to get a scrambled message of length nb bits is given by:  $ER = nbee$

### 6.1. Calculation FOR DYNAMIC REDUNDANCY

Administration OF MULTIPATH ROUTING The target of element repetition administration is to progressively distinguish and apply the best repetition level in terms of way repetition (mp) and source excess (ms), and additionally the best interruption location settings as far as the number of voters (m) and the interruption summon interim (TIDS) to augment MTTF, because nature's domain progressions to include parameters including SN/CH hub thickness ( $\lambda_{sn}/\lambda_{ch}$ ), radio extent (rsn/rch), and SN/CH catch rate ( $\lambda_c$ ). Our calculation for element repetition administration of multipath steering is circulated in nature. Figs. 2 and 3 depict the CH and SN execution conventions, individually, for overseeing multipath steering for interruption tolerance to augment the framework lifetime. They detail control activities taken by individual Sns and Chs because of powerfully evolving situation

All hubs in the framework demonstration occasionally to a "TD clock" occasion to alter the ideal parameter setting in light of evolving situations. . The ideal outline settings as far as ideal TIDS, m, ms, and mp are decided at static outline time and prestored in a table over detectable scopes of information parameter values. As there is no base station in the framework, the obligation of performing a table lookup operation with insertion and/or extrapolation methods connected to focus the ideal outline parameter settings will be accepted by Chs. The activity performed by a CH upon a TD clock occasion incorporates (a) changing CH radio reach to keep up CH network (line 4 in Fig. 2); (b) deciding TIDS, m, ms, and mp (line 5 in Fig. 2) based on the sensed natural conditions at runtime; and (c) informing Sns inside the bunch of the new TIDS and m settings. The activity performed by a SN upon this TD clock occasion is to change its radio extent to keep up SN integration inside a group . The move made by a SN after accepting the control bundle from its CH is to upgrade the new TIDS and m settings for interruption identification. At the point when the TIDS clock occasion happens, every hub in the framework utilizes its current TIDS and m settings to perform interruption identification. The TIDS clock occasion and the move made are pointed out for a Sensor Network

## 7. Execution Assessment

In this area, we exhibit numerical information got as a consequence of applying Eq. (1). Table II records the set of information parameter qualities describing a grouped HWSN. Our illustration HWSN comprises of 3000 SN hubs and 100 CH hubs, conveyed in a square zone of A2 (200m × 200m). Hubs are conveyed in the zone after a Poisson process with thickness  $\lambda_{sn} = 30$  hubs/(20 × 20 m<sup>2</sup>) and  $\lambda_{ch} = 1$  hub/(20×20 m<sup>2</sup>) at arrangement time. The radio extents rsn also rch are rapidly balanced between 5m to 25m and 25m to 120m individually to keep up system network. The starting vitality levels of SN and CH hubs are ESN 0 = 0.8 Joules and ECH 0 = 10 Joules so they deplete vitality at about the same

## 8. CONCLUSION

In this paper we performed a tradeoff investigation of vitality utilization vs. Qos pick up in unwavering quality, convenience, and security for excess administration of grouped heterogeneous remote sensor systems using multipath directing to reply client inquiries. We created a novel likelihood model to dissect the best excess level regarding way repetition (mp) and source excess (ms), and in addition the best interruption discovery settings regarding the amount of voters (m) and the interruption summon interim (TIDS) under which the lifetime of a heterogeneous remote sensor system is amplified while fulfilling the dependability, auspiciousness and security prerequisites of inquiry transforming applications in the vicinity of problematic remote correspondence and malevolent hubs. At long last, we connected our dissection results to the outline of an element excess administration calculation to recognize and apply the best outline parameter settings at runtime accordingly to environment progressions to draw out the framework lifetime

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