Hybrid Scheduling Algorithms Using Round Robin And Short Job First

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ABSTRACT

Scheduling is considered one of the most important topics in operating system. Generally, it helps to choose certain processes for execution. This paper will present a new integrated scheduling algorithm for multiprocessor system using Round Robin and Shortest Job First algorithms. Based on this integration, the result will be presented and compared with the previous algorithms. The new integrated scheduling algorithm was found to be simple and significant by reducing the waiting time. The result shows an improvement in this algorithm up 3.4% in terms of waiting time .

Key words Scheduling algorithms, hybrid algorithms, Round Robin(RR), Shortest Job First (SJF).

1. INTRODUCTION

Operating system is an interface between end user and system hardware, modern operating system become more complex than legacy operating system, especially when the turn of a single task into multiple tasks. CPU scheduling it is one of the most important works by operating system and the process time which is very important and also help in the selection process by which disease execution of several processes exist within the whole queue waiting for their turn to primitive execution.

The Central Processor Unit (CPU) executes the processes within the system operations one after the other, and dividing the time between these operations and serve each Process as necessary [5]. The scheduling algorithms have different properties, and may be one of these algorithms better than the other, the final goal in scheduling is to obtain optimum scheduling. In Single-processor computer systems, the efficiency of the scheduling algorithms rely on fairness, CPU Utilization, CPU Throughput, Turnaround Time, Waiting Time and Response Time [4]. Scheduling tasks on processors is considered one of the most important issues studied to make processors operate without being idle for long time. This increased the interest in studying scheduling and its

algorithms [6] .[2] stated that scheduling policy is considered one of the main factors that affect the performance. It helps to choose which task should be selected first from ready queue to run. Round Robin (RR) scheduling algorithm is widely used furthermore, its execution exceptionally Reliability on a Quantum estimate Qt, which is a predefined measure of time appointed by CPU to each errand to be executed. Notwithstanding, the execution debases regarding a normal holding up time (A WT), a normal Pivot time (ATT) and various Setting Switches (NCS). [2] presented new enhanced dynamic Round Robin booking calculation to lessen the normal holding up time, turnaround time and the quantity of setting changes with a specific end goal to enhance the framework general execution. It also talented a comparative analysis between several existing Round Robin algorithms based on the average time for waiting and turn-around and number of context switches. [1] studied the static scheduling issue for the independent tasks on a homogenous multiprocessor system. They develop an algorithm based on Bees Colony Optimization. They compare their algorithm with a previous one inspired also by the bees behavior for the same purpose. Their algorithm shown that the imposed algorithm gives the best solution for the scheduling problem, in most cases, and improves the traditional BCO algorithm.

2. PROBLEM OF STUDY

Task Scheduler is considered as an important issue addressed because All work carried out by processors seeking to reduce the time of completion. Most scheduling algorithms are having trouble finding the optimal time to perform tasks through Distribution blindness processors, as well as the cost consumed by the algorithm to find the optimal solution.

This paper presents a new hybrid algorithm that mix the Round Robin scheduling algorithm with the Short job first Planning algorithm to decrease the Normal Holding up time, turnaround time and the number from claiming setting switches so as on move forward the framework general execution. We will compare between the existing Round Robin algorithms and our proposed hybrid algorithm based on the average time.

3. THE PROPOSED ALGORITHM

This section will introduce the proposed hybrid algorithm; the algorithm suggests the merge of two algorithms (Round-Robin, Short job first); we enter the smallest value (the needed execution times for processes) rather than the distribution of the quantum on all operations without considering the value of time needed by each process.

4. THE PROPOSED ALGORITHM STEPS

1. we assume the arrival time of a set of processes to be at the same time.

2. We identify the value of Q (quantum), which means amount of execution of processes evenly.

3. the algorithm rank the process based on their CPU Burst Time (Ascending); thin it takes the smaller value between CPU Burst Times for the processes, then grant it the quantum value, if it is finished then it will move to the next process, if not finished within the quantum granted to it; then it will wait for the next loop. We move to the next process.

Experiment 1:

Process name	ВТ
Pr1	22ms
Pr2	18ms
Pr3	9ms
Pr4	10ms
Pr5	5ms

According to simple Round Robin scheduling: Round Robin quantum =5 Gantt chart:

| Pr |
|----|----|----|----|----|----|----|----|
| 1 | 2 | 3 | 4 | 5 | 1 | 2 | 3 |

Average waiting time=34 ms

According to The proposed algorithm scheduling: Round Robin quantum =5 Gantt chart.

Ount	it onui	ι.									
Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	PrAccontingPto simple Ro
5	3	4	2	1	3	4	2	1	2	1	2 Round Robin quantum

Pr

4

Pr

1

Pr 2

Average waiting time=25 ms

The method used	Average waiting time
Round Robin	34 ms
The proposed algorithm	25 ms

Experiment 2:

Process Id	AT	BT
Pr1	0ms	80ms
Pr2	0ms	45ms
Pr3	0ms	62ms
Pr4	0ms	34ms
Pr5	0ms	78ms

According to simple Round Robin scheduling:

Round Robin quantum =25 Gantt chart.

Gan	u chai	ι.									
Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr
1	2	3	4	5	1	2	3	4	5	1	3

Average waiting time=187.2 ms

According to The proposed algorithm scheduling:

Round Robin quantum =25

Gantt chart.

0		••									
Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr	Pr
4	2	3	5	1	4	2	3	5	1	3	5
		• •		1 - 1							

Average waiting time = 164.6 ms

The method used	Average waiting time
Round Robin	187.2 ms
The proposed algorithm	164.6 ms

Experiment 3:

Pro	ocess i	name		AT	BT				
Pr1				0ms	77ms				
Pr2	2			Oms	54ms				
 Pr	Pr	Pr							
l Pr3	2	1		0ms	45ms				
Pr4	ļ			0ms	19ms				
Pr5				0ms	14ms				

ound Robin scheduling: n =25

Gantt chart:

P1	P2	P3	P4	P5	P1	P2	P3	P1	P2	P1

Average waiting time = 117.4 ms

According to The proposed algorithm scheduling: Round Robin quantum =25

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	nd Ro					117.4					Round Robin quantum =25												
The	propo	sed alg	orith	n		71.4	ms				Ga	ntt ch							_		-	-	
Expe	erimer	nt 4 :									Pr 1	Pr 2	Pr 3	4		Pr 5	Pr 1	Pr 2	Pr 3	Pr 4	Pr 5	Pr 1	Pr 2
Proc	ess na	ame	A	T		I	BT				Ac	erage cordin und F	g to T	The p	rop	osed a		hm sc	heduli	ng:			
Pr1			0	ms		1	05ms				Ga	ntt ch	art:				D	<u> </u>			D	D	Б
Pr2			0	ms		9	0ms				Pr 3	Pr 4	Pr 2	P 5		Pr 1	Pr 3	Pr 4	Pr 2	Pr 5	Pr 1	Pr 2	Pr 5
D.2							0				Av	erage	waitii	ng tir	ne =	166	ms.		·	•	•	•	_1
Pr3			0	ms		e	0ms				Th	e met	hod u	sed			A	verag	e waiti	ing tin	ne		
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Pr5			0	ms		3	5ms				Ex	perim	ent 6	:									
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		bin qu	antun	n =25	•			•			Pr4	ł			Or	ns			10ms				
	tt char			1	1			1	1		Pr	ī —		1	Or	ne		1	3ms				
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5	4	3	2	1	5	4	3	2	1	3	Ac	cordin	g ^ź to s	impl	e Ro	ound	Robir	sche	duling	:			
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Pr1			0	ms		9	92ms					und F			;th			9.2 ms					
Pr2			0	ms			70ms				In	e prop	osea a	ugor	unn	1	23	3.2 ms					
Pr3			0	ms			35ms												een the			m	
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Experimental results

A table showing the difference between the results :

No. of experiment	Round Robin	The proposed
		algorithm
Experiment 1	34 ms	25 ms
Experiment 2	187.2 ms	164.6 ms
Experiment 3	117.4 ms	71.4 ms
Experiment 4	214 ms	167 ms
Experiment 5	173.4 ms	166 ms
Experiment 6	29.2 ms	23.2 ms

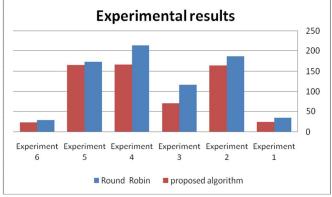


Figure 1: Experimental all results

5. 2.1 PSEUDO CODE THE PROPOSED ALGORITHM

```
burst time[n] \leftarrow 0
arrival time \leftarrow 0
numprocess[n] \leftarrow 0
turn[n] \leftarrow 0
wait[n] \leftarrow 0
temp \leftarrow 0
current time \leftarrow 0
tslice \leftarrow 0
waittime = 0
turntime = 0
avgwaititme=0.0
avgturntime=0.0
Read bursttime[n]
For i \leftarrow 0 to n-1
For j ← I to n
temp \leftarrow burst time[i]
```

burst time[i] ← burst time[j] burst time[j] ← temp end for For $i \leftarrow 0$ to n-1For j \leftarrow I to n Order burst time[i] temp ← burst time[i] burst time[i] \leftarrow burst time[i] burst time[j] ← temp end for j if bursttime[i] != 0 if tslice < bursttime[i] temp \leftarrow bursttime[i] < bursttime[j]) bursttime[i ← bursttime[i]-tslice[i] currenttime \leftarrow currenttime+tslice else arrivaltime \leftarrow arrivaltime+bursttime[i]; currenttime \leftarrow currenttime+ bursttime[i] bursttime[i ← bursttime[i]- bursttime[i] end else if bursttime[i] == 0turn[i] ← currenttime end if end if else currenttime++ end else for $i \leftarrow 0$ to n wait[i] ← turn[i]-burst[i]-arrivaltime $turn[i] \leftarrow turn[i]$ -arrivaltime waittime \leftarrow waittime+wait[i] turnrime ← turntime+turn[i] end of for avgwaittime \leftarrow waittime/n avgturntime ← turntime/n

5. CONCLUSION AND FUTURE WORK

The study found an efficient algorithm to schedule a series of operations on Optional number of independent processors, and can be applied in the case of this algorithm Processes independent, as well as whether the execution times equal Operations.

The proposed variant of RR algorithm drastically decreases context switching . The proposed algorithm performs better than RR scheduling algorithm with respect to average waiting time . our proposed algorithm can be further in investigated to be useful in providing more and more task-oriented results in future along with developing adaptive algorithms to fit the varying situations in today multifaceted complex working of operating system.

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